

EECS 755 - Final Exam

Spring Semester 2024

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Exercise 1 Which of the following statements are true? (1pt each)

1. The *desruct* tactic implements proof-by-cases.
2. The *Inductive* construct defines a set of all elements of a type.
3. The *Definition* construct can define a recursive function.
4. A nonterminating tactical introduces an inconsistency in Coq.
5. A tactical takes tactics as arguments.
6. $t1 ; t2$ is a tactical that applies $t1$ and subsequently applies $t2$ to the first goal resulting from $t1$.
7. $\text{try } t$ tries tactic t and if it fails returns the proof goal to its previous state.
8. The *inversion* implements backwards reasoning while *apply* implements forward reasoning.
9. *True* is a proposition with exactly one proof.
10. $\text{not } A$ unfolds to $(\text{not } A) \rightarrow \text{False}$
11. *split* takes a goal of the form $A \leftrightarrow B$ and produces two separate goals $A \rightarrow B$ and $B \rightarrow A$.
12. In the Hoare triple $\{\{Q\}\}c\{\{P\}\}$ specifies that Q must be true before c executes for P to hold afterwards.
13. *state* is a mapping from variable names to their values.
14. In *The Adventures of Buckaroo Bonzai John Lithgow* plays a mad scientist trying to return to the 8th dimension using a device called an overthruster.

Exercise 2 In this exercise we will be thinking command definition. You will define a new *do-until* loop that executes its body and then checks its termination condition. Unlike *while* which checks before. You will also be thinking about specifying read-only memory.

1. Extend the inductive proposition definition of IMP to define *do c until b*
2. Can you extend the *eval* function for IMP to include *do c until b*? If so, do it. If not explain why not.

3. What is the difference between IMP com expressions and the `aexp` and `bexp` expressions? Specifically, what is the difference between how they are evaluated?
4. Most languages split their states to isolate read-only values from values that may change. Using Coq, specify a new `state` that does this.
5. How would we prove that an IMP program does not attempt to modify its read-only values?

Exercise 3 In this exercise we're going to think a bit more about the `until` loop and equivalence with `while`. Also looking at proof commands for reasoning about language constructions.

1. Using `while` define a construct that is equivalent to `until` from the previous problem.
2. What theorem should you prove to show this new `until` and your `while` construction are equivalent.
3. When reasoning about some program of the form `c1 ; c2` in the assumptions of a proof, why might `inversion` or `be good options`? Are we reasoning forwards or backwards?
4. When reasoning about some program of the form `c1 ; c2` as the goal of a proof, why would `eapply E_Seq` be a good option? Are we reasoning forwards or backwards?
5. Assume that we've defined a new, swanky optimizer for IMP called `swanky` that inputs an IMP program and returns an optimized IMP program. What theorem would you prove to show the optimizer is correct?

Exercise 4 Finally we're going to think a bit about Hoare Logic. We're going to push a bit here and think about concurrency and what that might mean.

1. Briefly explain what the notation $\{\{X=m\}\}X:=X+1\{\{X=m+1\}\}$ means?
2. If we have some command defined by $\{\{True\}\}c\{\{Y=n \text{ AND } X=m\}\}$, can we evaluate $c;X:=X+1$? Why or why not? (Your answer need not be formal.)
3. Let's define a new command `c1 || c2` that behaves like sequence except that `c1` and `c2` execute in parallel. `c1` and `c2` start in the same state and should end in a state that reflects execution of both commands. Define an inference rule for this command and a Hoare Logic rule.
4. Let's say that `c1 || c2` is serializable if its execution result is the same no matter who goes first. Can you capture serializability using Hoare Logic? Specifically, how would you prove that `c1 || c2` is serializable?

5. Let's say that $c1$ and $c2$ interfere with each other when the state resulting from $c1 \parallel c2$ is inconsistent. Can you define interference using Hoare Logic? Specifically, how would you prove that $c1 \parallel c2$ exhibits interference?