

EECS 755 - Final Exam

Fall Semester 2025

December 5, 2025

Exercise 1 Circle the numbers associated with statements in the following list that are true. (1pt each)

1. The **destruct** tactic implements proof-by-cases.
2. The **Inductive** construct defines the set of all elements of a type.
3. The **Definition** construct can define a recursive function.
4. A non-terminating tactical introduces an inconsistency in Coq.
5. A tactical takes tactics as arguments.
6. $t_1; t_2$ is a tactical that applies t_1 and applies t_2 to the first goal resulting from t_1 .
7. `try t` attempts tactic t and if it fails returns the proof goal to its previous state.
8. `apply L` in H where H is a hypothesis implements forward reasoning while `apply L` implements backward reasoning
9. `not A` unfolds to $A \rightarrow \text{False}$
10. `split` takes a goal of the form $A \leftrightarrow B$ and produces two separate goals $A \rightarrow B$ and $B \rightarrow A$.
11. In *The Adventures of Buckaroo Bonzai* John Lithgow plays a mad scientist trying to return to the 8th dimension using a device called an overthruster.

Exercise 2 *Using an Inductive construction we will define a simple stack data structure. (3pts each)*

1. *Using Inductive define a stack that can contain any single type. You should minimally define an empty stack and a push operation.*
2. *Define top and pop over your new data structure.*
3. *How would you prove that the top of a stack resulting from pushing x is always x ?*
4. *How would you prove that pushing never results in an empty stack?*
5. *What is the difference between the definition of **stack** and the definition of **nat**? Specifically, could you use one to implement the other?*

Exercise 3 Assume the following AST for a simple language of Booleans and if: (3pts each)

```
Inductive L : Type :=
| ttrue : L
| tfalse : L
| tand : L -> L -> L
| tnot : L -> L
| tif : L -> L -> L -> L.
```

1. Define an evaluation function, *teval*, for *L* that produces a value of type *bool*.

2. Define an evaluation relation, *tevalR*, for *L* that relates an expression to a value of type *bool*.

3. Assuming *tif* is a traditional if-expression, write a relation that describes an optimization that replaces *(tif l1 l2 l2)* with *l2*.

4. Assuming *tif* is a traditional if-expression, write an optimization function that replaces *(tif l1 l2 l2)* with *l2*.

5. What theorem would you prove to show the optimization relation is correct?

6. What theorem would you prove to show the optimization function is correct assuming the optimization relation is correct?
7. We want to prove $(\text{forall } t1 \ t2:L, (\text{teval } (\text{tif } t1 \ t2 \ t2)) = (\text{teval } t2))$. Should we start with induction or destruct or something else? (Not intros.)
8. We want to prove $(\text{forall } t1, \text{teval } (\text{tand } t1 \ \text{tfalse})) = \text{false}$. Should we start with induction or destruct or something else? (Not intros.)
9. We want to prove $(\text{forall } t1, \text{teval } (\text{tand } t1 \ \text{tfalse})) = \text{false}$. Now tell me what intros does to the theorem.
10. We have now proven Theorem T: $(\text{forall } t1, \text{teval } (\text{tand } t1 \ \text{tfalse})) = \text{false}$. Can we use the resulting theorem T with either apply or rewrite?

Exercise 4 The definition of IMP from class defines state as a *Map* of variable names to their values with functions to update state and retrieve values from state. Evaluators for IMP accept and produce states. Answer the following about IMP and state: (3pts each)

1. When reasoning about some program with *c1* ; *c2* in the assumptions of the proof, would using **destruct** be a good option? If so, are we reasoning forwards or backwards?

2. When reasoning about some program with *c1* ; *c2* in the goal of the proof, would using **apply E_Seq** be a good option? If so, are we reasoning forwards or backwards?

3. Assume that we've defined a new, swanky optimizer for IMP called **swanky** that inputs an IMP program and returns an optimized IMP program. What theorem would you prove to show the optimizer is correct?

4. Our definition of IMP defined three sub-languages - *aexp*, *bexp*, and *com*. The interpreters for *aexp* and *bexp* accept a term and produce a value. What does the interpreter for *com* accept and produce?

5. We say that a variable is **invariant** over program execution if it's value does not change. If we wanted to verify that a variables *X* and *Y* are invariant during IMP program execution what would we prove? (You may assume the existence of *comR*, a relational definition for *com* evaluation.)